Optimisation



Derivative-free optimization: From Nelder-Mead to global methods

Definition

Optimizing a function is looking for the set of values of the variables that will maximize (or minimize) the function.

- Optimization is usually a very complex problem. There are many different techniques, each being adapted to a specific kind of problems.
- Here is no universal method, but a set of tools which requires a lot of experience to be used properly.

Optimization caracteristics

₭ Global/local optimization

- Global optimization is searching for the absolute extremum of the function over its entire definition domain
- Local optimization is looking for the extremum of the function in the vicinity of a given point

Stochastic / deterministic methods

- A stochastic method searches the definition domain of the function in a random way. Two succesive runs can give different résults.
- △ A deterministic method always walks the search space in the same way, and always gives the same results.

x²+**y**²























Part I Local deterministic methods

Derivation (deterministic)

₩When it is possible to compute and solve f'(x)=0, then we know that the extrema of the function are in the set of solutions.

His method can only be used for very simple analytic functions

Gradient method (deterministic and local)

If f(X) is a real valued function of a real valued vector X, and we can calculate f'(X), we compute:

 $X_{n+1} = X_n - a f'(X_n), a>0$ The best choice of a>0 is done by minimizing:

 $G(a) = f(X_n - a f'(X_n))$ It's usually impossible to solve the above equation and approximate methods are used.



Local, deterministic, order 2, methods.

To accelerate computation we use the computation of the first and second order derivatives of the function

How we want to be able to compute both, in a reasonnable amount of time.

Local, deterministic, order 2, methods.

 $\begin{aligned} & \texttt{H}f(y) = f(x) + f'(x) (y-x) + \frac{1}{2} f''(x) (y-x)^2 + d \\ & \texttt{W}e \text{ minimize the y quadratic form:} \\ & \bigtriangleup f'(x) + f''(x)(y-x) = 0 = > y = x - \frac{f'(x)}{f''(x)} \end{aligned}$

$$x_{n+1} = x_n - f'(x_n) / f''(x_n)$$

Known as Newton method

Convergence is (much) faster than the simple gradient method.

Newton



Deterministic method: BFGS

#BFGS approximates the hessian matrix without explicitly computing the hessian

It only requires knowledge of the first order derivative.

It's faster than gradient, slower (but much more practical) than Newton

#One of the most used method.

BFGS



Local deterministic: Nelder-Mead simplex

Works by building an n+1 points polytope for an n variables function, and by shrinking, expanding and moving the polytope.

Here's no need to compute the first or second order derivative, or even to know the analytic form of f(x), which makes NMS very easy to use.

#The algorithm is very simple.

Nelder-Mead simplex

#Choose n+1 points $(x_1, ..., x_{n+1})$ **#**Sort: $f(x_1) < f(x_2) \dots < f(x_{n+1})$ **#**Compute barycenter: $x_0 = (x_1 + ... + x_n)/n$ \Re Reflection of x_{n+1}/x_0 : $x_r = x_0 + (x_0 - x_{n+1})$ $\Re \inf f(x_r) < f(x_1), x_p = x_0 + 2(x_0 - x_{n+1}).$ If $f(x_p) < f(x_r),$ $x_{n+1} < -x_{e}$, else $x_{n+1} < -x_r$, back to sort. $\Re If f(x_n) < f(x_r), x_c = x_{n+1} + (x_0 - x_{n+1})/2. If f(x_c) < f(x_r)$ $x_{n+1} < -x_{c}$, back to sort **Herwise:** $x_i < -x_0 + (x_i - x_1)/2$. Back to sort.

Nelder Mead



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Part II Global stochastic methods

Stochastic optimization

Bo not require any regularity (functions do no even need to be continuous)

- **#**Usually expensive regarding computation time, and do not guarantee optimality
- Here are some theoretical convergence results, but they usually don't apply in day to day problems.

Simulated annealing

Generate one random starting point x₀inside the search space.

Build $x_{n+1} = x_n + B(0,s)$ Compute: $t_{n+1} = H(t_n)$ ∴ If $f(x_{n+1}) < f(x_n)$ then keep x_{n+1} ∴ If $f(x_{n+1}) > f(x_n)$ then : ∴ If $|f(x_{n+1}) - f(x_n)| < e^{-kt}$ then keep x_{n+1} ∴ Si $|f(x_{n+1}) - f(x_n)| > e^{-kt}$ then keep x_n

Important parameters

#H (the annealing schedule):

- Too fast=>the algorithm converges very quickly to a local minimum
- ☐ Too slow=>the algorithm converges painfully slowly.
- Deplacement: B(0,s) must search the whole space, and mustn't jump too far or too close either

Efficiency

SA can be useful on problems too difficult for « classical methods »

#Genetic algorithms are usually more
efficient when it is possible to build a
« meaningful » crossover

Genetic algorithms (GA)

Search heuristic that « mimics » the process of natural evolution:

Reproduction/selection

Crossover

Mutation

KJohn Holland (1960/1970)

#David Goldberg (1980/1990).

Coding / population generation

#If x is a variable of f(x), to optimize on the interval [x_{min},x_{max}].

 $\text{HWe rewrite x :} 2^{n} (x - x_{\min}) / (x_{\max} - x_{\min})$

#This gives an n bits string:

✓For n=8: 01001110

✓ For n=16: 0100010111010010

∺A complete population of N (n bits string) is generated.



#Two parents : △01100111 △10010111 \Re One crossover point (3): △011 00111 △100 10111 **H**Two children: △011|10111 △100 00111

Mutation

Reproduction/selection

- **#**For each x_i compute $f(x_i)$ **#**Compute $S=\Sigma(f(x_i))$
- **H**Then for each x_i :

 $\square p(x_i) = f(x_i)/S$

#The n elements of the new population are picked from the pool of the n elements of the old population with a bias equal to $p(x_i)$.

#Better adapted elements are more reproduced

Exemple de reproduction

%#f(x)=4x(1-x)%#</pre

Séquence	Valeur	U(x)	% de chance	% cumulés	Après
_			de reproduction		reproduction
10111010	0.7265625	0.794678	0.794678 / 2.595947 = 0.31	0.31	11011110
11011110	0.8671875	0.460693	0.460693 / 2.595947 = 0.18	0.31+0.18=0.49	10111010
00011010	0.1015625	0.364990	0.364990 / 2.595947 = 0.14	0.49+0.14=0.63	01101100
01101100	0.4218750	0.975586	0.975586 / 2.595947 = 0.37	0.62+0.37=1.00	01101100
=		2.595947			

AG main steps

Step 1: reproduction/selection

- Step 2: crossing
- Step 3: mutation
- Step 4: End test.


#Fact: in the « simple » AG, the fitness of an element x is equal to f(x)

Instead of using f(x) as fitness, f is « scaled » by using an increasing function.
Exemples:

 \bigtriangleup 5 (f(x)-10)/3: increase selection pressure

 \bigcirc 0.2 f + 20 : diminishes selection pressure

#There are also non-linear scaling functions

Scaling examples





Selection pressure can induce too fast convergence to local extrema.

Sharing modifies fitness depending on the number of neighbours of an element:

$$\triangle f_s(x_i) = f(x_i) / \Sigma_j s(d(x_i, x_j))$$

Ad(x_i,x_j) is a distance measurement between i et j



How the second secon

#General shape of s:



Bit string coding problem

How wery different bit strings can represent elements which are very close to each other:

☐ If encoding real values in [0,1] with 8 bits:

≥ 10000000 et 01111111 represent almost the same value (1/2) but their hamming distance is maximal (8).

△ Necessity to use Grey encoding.

Using a proper coding

#For real variable functions, real variable
encoding is used

#Crossover:

 $\Delta \alpha$ randomly picked in [0.5,1.5]

#Mutation:

$$Y_1 = X_1 + B(0,\sigma)$$



Aircraft conflit resolution



Modeling

Conly one manoeuver maximum by aircraft
10, 20 or 30 degrees deviation right or left
Then return to destination
Offset
Variables: 3 n

- △T0: start of manoeuver
- △T1: end of manoeuver
- △A: angle of deviation
- **#**Uncertainties on speeds











0: 0: 6 - Prevision 20 mn - Opti toutes les 4 mn EN COURS - Incertitude 5%

0.383436



Results

Résultats

- Résout des gros conflits (30 avions)
- Intégration dans un outil de simulation (CATS/OPAS)
- Testé sur des journées de trafic réel
- Peu de restrictions sur la modélisation
- Pas de garantie d'optimalité



Traveling Salesman Problem (TSP)



TSP: crossover



TSP: new crossover



TSP: mutation



Ant Colony Optimization (ACO)

Mimic the ants trying to find the shortest path to food



ACO

#Ants deposit pheromones according to the quality of path

- Ants more likely to follow paths with the most pheromones
- Evaporation process to prevent early convergence
- Stop when no more improvement

ACO for the TSP

#Each ant builds a path **#**Choice of next city influenced by pheromones already present **#**Ants deposit pheromons on the path chosen \Re At each iteration, pheromons evaporate



Differential Evolution

[₩]Pick:

△NP vector elements population with n variables

→ F in [0,2] (differential weight)

△CR in [0,1] (Crossover probability)

#For each vector element x

 \boxtimes Pick randomly 3 distinct vectors a,b,c in population \boxtimes Pick a random index R in [1,n]

Kertein Karakan Kar

• If $r_i < CR$ or i=R then $y_i=a_i+F(b_i-c_i)$ else $y_i=x_i$

 \boxtimes If f(y) better than f(x) replace x by y in population

Other evolutionary techniques

%Particle Swarm optimization
%Evolutionary strategies
%Genetic Programming
%



Part III Global deterministic methods

B&B and interval programming (global deterministic methods)

₩With:

- $\begin{aligned} f(x,y) &= 333.75 \ y^6 + \ x^2 \ (11 \ x^2 y^2 \ y^6 \ 121 \ y^4 \ 2) \\ &+ \ 5.5 \ y^8 + \ x \ / \ (2y) \end{aligned}$
- ₭ If we compute f(77617,33096), we get 1.172603.
- **#**The correct value is -0.827396.
- Interval program was initially designed to circumvent improper rounding.

Elementary operations

- ₭ If X=[a,b] and Y=[c,d]
- \Re X+Y=[a+c,b+d] and X-Y=[a-d,b-c]

<mark>₩</mark> X*Y=

- [ac,bd] si a>0 et c>0
- [bc,bd] si a>0 et c<0<d</p>
- [bc,ad] si a>0 et d<0</p>
- [ad,bc] si a<0<b et c>0
- □ [bd,ad] si a<0<b et d<0</p>
- [ad,bc] si b<0 et c>0
- [ad,ac] si b<0 et c<0<d</p>
- [bd,ac] si b<0 et d<0</p>
- [min(bc,ad),max(ac,bd)] si a<0<b et c<0<d</pre>

Divide

#R is extended using $+\infty/-\infty$ **#**X/Y=

- $[b/c, +\infty]$ if b<0 and d=0
- \square [- ∞ ,b/d] and [b/c,+ ∞] if b<0 and c<0<d
- \square [- ∞ , + ∞] if a<0<b
- \square [- ∞ ,a/c] if a>0 et d=0
- \square [- ∞ ,a/c] and [a/d, + ∞] if a>0 et c<0<d

 $[a/d, +\infty]$ if a>0 and c=0

Other operations

#All operations can be extended to interval arithmetic.

#For monotonous functions:

 \square F([a,b])=[f(a),f(b)] if f is increasing

 $\mathbb{P}F([a,b]) = [f(b),f(a)]$ if f is decreasing

 \triangle Example: Exp([a,b])=[e^a,e^b]

Composing functions is done by composing interval extensions of these functions

Problems

 \Re If X=[a,b], X-X = [a-b,b-a]<>[0,0]! High In the same way (X-1)(X+1) <> X²-1 $\Re([0,2]-1)([0,2]+1)=[-1,1]*[1,3]=[-3,3]$ $\mathbb{H}[0,2]^2-1=[0,4]-1=[-1,3]$ **#**Associativity is preserved: $\triangle A + (B+C) = (A+B) + C$ $\triangle A(BC) = (AB)C$

Branch and bound

Generic name for all methods that divide and cut part of the search space.

Here, search space is divided by cutting intervals in two, and bounds are generated by estimating the function value over each sub-interval.

Minimization

Set: L<-{[a,b]} et e<-estimator of f on [a,b]

Extract I=[c,d] top of L. If e<c, redo. If I is too small, redo. If L is empty: end.

- △Build $I_1 = [c, (c+d)/2]$ and $I_2 = [(c+d)/2, d]$.
- Compute $F(I_1) = [x_1, y_1]$, $F(I_2) = [x_2, y_2]$, e_1 et e_2 .
- Set e=min(e,e1,e2)
- \square If $x_1 < e$ then insert I_1 in L
- \square If x₂<e then insert I₂ in L

Back to start.

Computation of the estimator

- △Sampling: take n points equally spaced in X
- Stochastic: draw randomly n points in X
- Computer f'(x) and F'(X) et check if the sign of f'(x) is the same on X = > f is monotonous and the extremum is on one side of the interval

How to sort the list of intervals

Many ways:
First In First Out
Largest first
Best estimator first
Smaller lower bound first
etc...

End test

₭ Many ways:

△The size of the interval is smaller than a defined value

△The size of the image of the function is smaller than a defined value

Etc...
More than one dimension

For a multiple dimension functions, cutting is done on each variable in turn.
It's usually the largest interval which is cut first.

#The end test is modified accordingly.



When to use it

Here program computing the function can be « easily » extended to interval arithmetic.

∺Method efficient when there are not too many variables.

\mathbb{H} In theory, computation time grows as $2^{\mathbb{N}}$ with N being the number of variables.



Part IV Cooperation

Cooperative algorithm

#IBBA thread

☐Gets from shared memory best EA element

 \boxtimes =>speeds up the cutting process

△Sends to shared memory its best element

#EA thread

△Sends to shared memory its best element

Replace worst element with best IBBA element

∺Update thread

Updates admissible domains/cleans up IBBA queue
Projects EA elements into the closest box

Cooperative algorithm Griewank D=6



Cooperative algorithm Statistical results

	size	6	7	8	9	10
EA	Found	100	94	92	83	15
	Mean (sec)	204	864	972	1340	1678
	Sigma (sec)	92	356	389	430	34
IBBA	Found	71	0	0	0	0
	Mean (sec)	284				
	Sigma (sec)	192				
Cooperative	Found	100	100	100	100	100
	Mean (sec)	50	62	156	215	267
	Sigma (sec)	18	47	85	317	105

Cooperative algorithm

Seful when the extremum has to be proved

Same constraints as the IBBA

Needs code that can be extended to interval arithmetics